# GNS: Abstract Syntax for Natural Languages 

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Summary. This paper presents an overview of General Natural Syntax (GNS), a formal theory of general explicative power that generalizes and formalizes syntactic concepts in order to offer a general notion of syntax that is independent of any particular language.

## 1 General Natural Syntax: Basic Idea

General Natural Syntax (GNS) is a formal theory of general explicative power that intends to formalize a general notion of syntax that is independent of any particular language. The idea behind this formalism is that it reduces syntactic constructions to a few principles related to their semantic functions, but which can be defined independently of semantics.

GNS generalizes and formalizes syntactic concepts so that they can be incorporated into a formal theory of general explicative power. As we have said, the general principles we proposed as elements of GNS are related to their
semantic function but definable independently of semantics. GNS representations stand between the concrete syntax of specific natural languages and the semantic function of syntactic phenomena. In this regard, GNS can be related to the so-called abstract syntax of programming languages [4]. Abstract syntax is an important tool in the translation process between programming languages: it provides a deep description of syntactic constructs and is independent of any particular syntactic encoding. The starting point for this approach was an algebraic notion of syntax, more abstract than a concrete syntax, placed at an intermediate level as a bridge between syntax and semantics. Symbols of an abstract syntax expression denote semantic functions, but at the same time, all the information is included so that the concrete syntax of the expression can be generated.

GNS is also related to the notion of tagged text, which is the basis of many markup languages (TeX, HTML, XML). Also here, abstract expressions are used that mix the pure textual information with information about text format and visualization (fonts, dimensions, paragraphs, etc.). So, the structure of a text is represented in a way that in many aspects is independent from the way a text is realized by a compiler or a browser. Marked texts abstractly express properties and relations of the textual units at a logical level.

Considering all the above ideas and taking into account that the use of algebraic formalization has a strong and deep tradition in natural language analysis within the fields of semantics [15, 14] and morphology [13], what we propose here is to develop an analogous logic-algebraic approach at an intermediate level between semantics and morphology: the level of the abstract or general syntax.

## 2 Formal Prerequisites

In this section we provide the formal prerequisites that are needed to understand the formalization presented in this paper. For further information on the theory of formal languages and mathematical logic see [17, 16, 2].

A nonempty set $V$ of symbols or letters is called an alphabet. A word or a string over an alphabet $V$ is an element of the free monoid $V^{*}$ generated by the symbols of $V$ under a binary associative operation of concatenation (denoted by juxtaposition). The empty word $\lambda$ is the neuter element with respect to concatenation $(x \lambda=\lambda x=x)$. The length of a string $x \in V^{*}$ (the number of symbol occurrences in $x$ ) is denoted by $|x|$.

We use symbols $\rightarrow, \neg, \wedge, \vee, \leftrightarrow, \forall, \exists, \equiv,=$ with the standard syntactic and semantic first order logic meaning. We assume that the reader is familiar with logical notions such as variables, constants, predicates, functors, first order formulae and terms, and free and bound variables. A set of formulae, called axioms, and all their logical consequences constitute a theory (with respect to some notion of logical consequence or of logical deduction).

A model $\mathcal{M}$, or a (relational) structure, is given by: i) a set $D$, called domain of $\mathcal{M}$, ii) some elements $a, b, \ldots \in D$, called individual constants of $\mathcal{M}$, and iii) some operations $f, g, \ldots$ and relations $R, Q, \ldots$ over $D$ (an arity is associated with each operation and relation and it specifies its number of arguments). Usually $\mathcal{M}$ is indicated by:

$$
\mathcal{M}=(D, a, b, \ldots, f, g, \ldots, R, Q, \ldots)
$$

The set $\operatorname{Term}(\mathcal{M})$ of the terms over $\mathcal{M}$ is given by all the expressions that can be constructed, in the usual algebraic sense, by applying the operations of $\mathcal{M}$ to the individual constants of $\mathcal{M}$. For example, if $f$ has arity 1 and $g$ has arity 2 , then the following are terms over the model given above:

$$
f(a), \quad g(a, b), \quad f(g(a, b)), \quad g(a, f(a)), \quad f(g(a, f(a))), \quad \ldots
$$

An equation such as $g(a, f(a))=b$ means that by applying the operation $f$ to the constant $a$ we get an element of $D$, say $c$, and by applying $g$ to the pair $a$ and $c$ we get $b$. This means that $g(a, f(a))$ is considered as the denotation of the element of $D$ obtained by applying the operations following the way algebraic expressions are usually evaluated (in the order specified by parentheses). However, we can consider the term as denoting itself, i.e. a sequence of symbols for individual constants, operations, commas, and parentheses, disregarding any meaning of symbols. It is important to distinguish between these two aspects. If we want to be precise we write $g(a, f(a))$ to refer to the element of $D$ (if it exists) denoted by the term, while in the other case we write $\lceil g(a, f(a))\rceil$. However, in practice the context will indicate in which sense a term is used.

A string model or a monoidal model $\mathcal{M}$ is a model that has: i) a domain that includes a free monoid $V^{*}$ over an alphabet $V$, ii) constants that include a constant $\lambda$ for an empty string and for symbols of $V$, iii) operations that include a binary operation that on $V^{*}$ coincides with the string concatenation on $V^{*}$ (indicated by juxtaposition).

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A signature $\Sigma$ is essentially a set of symbols for denoting objects (individual constants), operations and relations of a relational structure (relations can be identified as operations that provide Boolean values as results). In $\Sigma$, any symbol of operation or relation is equipped with an arity that specifies the number of arguments of the corresponding operation or relation.

A string theory or a monoidal theory $\mathcal{T}$, over an alphabet $A$ and a signature $\Sigma$, is a first order logic theory of axioms $\Phi$ over the signature $\Sigma$ such that all the symbols of $A$ and the empty string $\lambda$ are individual constants of $\Sigma$. The concatenation operation of monoids is denoted in $\Sigma$, and $\Phi$ also includes the monoid axiom: $\forall x, y, z(x(y z)=(x y) z \wedge x \lambda=x \wedge \lambda x=x)$. In other words, a monoidal theory over an alphabet $A$ and a signature $\Sigma$ is a theory where terms include the strings over $A$. This means that concatenation is applied not only to the symbols of $A$, but to the terms that can be built on the signature $\Sigma$. This possibility provides syntactic constructs where not only symbols of an alphabet, but even terms, are concatenated.

A monoidal system or simply a monoidal indicates, generically, a string model or a string theory. Monoidals are good environments for defining the syntax of formal languages [11].

Let us indicate by $G N S(L)$ the terms, with their syntactic category, that can be constructed starting from the lexical items of $L$ in a suitable string theory that will be defined in the course of the paper. The alphabet of this theory is the set of usual Latin letters plus other special symbols that will appear in the axioms given in the next sections. Variables are indicated by letters $x, y, z, u, v, \ldots$ which may have apices or subscripts. Individual constants are indicated by special strings of capital letters. Operations and predicates are indicated by strings starting with a backslash $(\backslash)$. For a better reading, given a predicate $P$ we write $t: P(\mathrm{t}$ is of type $P)$, instead of $P(t)$, in fact, $t: P$ intuitively means that term $t$ belongs to the syntactic category $P$.

## 3 General Natural Syntax: the Formalism

GNS has been defined as a general formal framework given by some axioms according to which some operations are applied to strings of some categories and get other strings (of some specified categories). The formalism consists of:

- Eight basic categories;
- Thirty syntactic operations;

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- Forty axioms;
- A few hundred grammemes;
- A few thousand lexical items (basic lexicon).

In what follows we will present the above elements in order to provide an overview of GNS.

### 3.1 Three basic operations

In order to construct GNS, three basic operations have been assumed as the preliminary data of the analysis:

1. Conjugation ( $\backslash$ conjug): adds temporal and dynamic parameters to syntactic elements that can play verbal roles.
2. Determination ( $\backslash$ determ): adds spatial and contextual information to elements that can play nominal roles.
3. Predication $(\backslash p r e d)$ : is the basic sentence building construction.

We assume that we know when these operations can be applied to some arguments.

### 3.2 Eight basic categories

Taking into account the three basic operations introduced in the above section we consider the following eight basic categories:

1. Verb (Verb): We define the class Verb of elements $x$ such that, for some conjugative $y, \backslash \operatorname{conjug}(x, y)$ provides a result. Formally:

$$
\text { Verb }=\{x \mid \exists y z \backslash \operatorname{conjug}(x, y)=z, y \in \text { Conjugative }\}
$$

2. Noun (Noun): We define the class Noun of elements $x$, such that, for some determinative $y, \backslash \operatorname{determ}(x, y)$ provides a result. Formally:

$$
\text { Noun }=\{x \mid \exists y z \backslash \text { determ }(x, y)=z, y \in \text { Determinative }\}
$$

3. Substantive (Subst): A substantive is the result of the determination operation. Formally:

$$
\forall x y z(\backslash \operatorname{determ}(x, y)=z \wedge y: \text { Determinative } \rightarrow z: \text { Subst })
$$

4. Verbative (Verbt): A verbative is the result of the conjugation operation. Formally:

$$
\forall x y z(\backslash \operatorname{conjug}(x, y)=z \wedge y: \text { Conjugative } \rightarrow z: \text { Verbt })
$$

5. Proposition (Prop). A proposition is the result of a $\backslash p r e d$ operation that takes as arguments a substantive (Subst) and a verbative (Verbt) and provides a proposition as a result. Formally:

$$
\text { Prop }=\{z \mid \exists x y \backslash \operatorname{pred}(x, y)=z, x \in S u b s t, y \in V e r b t\}
$$

We write $x$ : Cat to state that $x$ is an expression of category Cat. We use Cat to indicate any of the previous categories. To those five basic categories we need to add the following three Ad categories (AdCat indicates any of them):
6. AdProp
7. AdNoun
8. AdVerb

It is important to note here that although the terms proposition, noun, verb, and substantive are taken from traditional grammatical, logical, and semantical analysis, the definition we provide of our basic types is completely formal, based on the assumption of some initial operations.

### 3.3 Grammemes

The eight basic categories above are those of full linguistic elements. To them we add a category of empty linguistic elements that we have called grammemes ${ }^{3}$. Grammemes determine the surface syntactic realization of the abstract syntax operations considered in GNS.

Syntactic operations defined in GNS allow us to construct complex linguistic expressions with one of the eithg basic categories described above. However, many syntactic operations need some additional parameters, which we have called grammemes. Grammemes are defined as elements that individuate the features required to evaluate terms of GNS(L) into expressions of L : that is, kinds of information that the grammar of $L$ requires to provide the

[^0]morphological realization corresponding to the abstract syntactic representation.

It is important to distinguish grammemes from particles. Particles are strings that are provided in the surface syntactic realization of the abstract syntax operations. Particles are strings that are inserted in a text as a consequence of applying a syntactic operation with some grammeme. Examples of grammemes are: singular (SING), plural (PLUR), present (PRES), presentprogressive (PPG), past (PAST), disjunctive (OR), definite (DEF), agent (AGENT), and first, second, third (I, II, III). The difference between grammemes and particles can be observed in the following English example: $\backslash \operatorname{conjug}$ (go, III, SING, PPG, PAST)= was going.

Grammemes are classified into several types. Although these types depend on the particular language we believe that they vary very little even within wide classes of languages. What is essential in the notion of grammemes is that they are a "closed set", which is a fixed and small number of elements not exceeding a few hundred units (apart from numeral and ordinal elements). Classes of grammemes in English, Italian, Spanish, and many other languages are listed in Table 1.

Note that some of the above classes consist of a sequence of grammemes, which depends on the language we are considering. For example, the class conjugative consists of a sequence which, depending on the language, may include types such as personal, number, gender, tense, modal-attitudinal, aspect....

### 3.4 Operations: Generalities

The syntactic operations we defined in GNS are related to fundamental semantic roles, but only the categories of their arguments and the categories of their results are essential in their determination. In this regard, a general syntactic construction is a kind of bridge between syntax and semantics. It does not deal with the particular morphological features of the final linguistic form of the expression that is the result of applying the syntactic operation, but at the same time, it does not concern the meaning of the final expressions and the way this meaning is related to those of the arguments of the operation. Therefore, a syntactic operation is defined by:
i) its name,
ii) the categories of its arguments, and
iii) the category of the resulting expression.

```
CLASSES OF GRAMMEMES
1. Conjugative:
    - Personal
    - Number
    - Tense
    - Mood
    - Gender
    - Aspect
    - Modal-Attitudinal
2. Determinative
3. Copulative
4. Circumstantial
5. Complementative
6. Locative
7. Cardinal
8. Ordinal
9. Deictic
10. Quantifier
11. Intensive
12. Comparative
13. Coordinative
14. Subordinative
```

Table 1. Classes of Grammemes

The syntactic operations defined can be classified as constructs of one of three general schemata:

1. Combination Schema puts together strings of different categories and provides a string of a category other than the categories of its components.
2. Expansion Schema takes as its argument a string of a given category and provides a bigger string of the same category.
3. Transcategorization Schema transforms a string of a given category into a string of another category.

Syntactic operations defined in GNS are listed in Table 2.

### 3.5 Axioms

$G N S$ is given by some axioms according to which operations are applied to strings of some categories and produce other strings (of some specified cat-

| N. | OPERATION | NOTATION |
| :---: | :---: | :---: |
| 1 | Conjugation | $\backslash$ conjug |
| 2 | Determination | $\backslash$ determ |
| 3 | Predication | $\backslash$ pred |
| 4 | Copulation | $\backslash$ copul |
| 5 | Modification | $\backslash$ modif |
| 6 | Apposition | $\backslash$ appos |
| 7 | Complementation | \ complem |
| 8 | Localization | $\backslash \mathrm{loc}$ |
| 9 | Coordination | $\backslash$ coord |
| 10 | Subordination | $\backslash$ subord |
| 11 | Negation | $\backslash$ neg |
| 12 | Passivation | $\backslash$ passiv |
| 13 | Adjectivation | $\backslash$ adj |
| 14 | Adverbialization | $\backslash$ adverb |
| 15 | Adproposition | $\backslash$ adprop |
| 16 | Relativization | $\backslash$ relat |
| 17 | Anacoluthon | $\backslash$ anacoluth |
| 18 | Numeralization | $\backslash$ num |
| 19 | Ordinalization | $\backslash$ ord |
| 20 | Deictification | $\backslash$ deixis |
| 21 | Anaphorization | $\backslash$ anaphor |
| 22 | Quantification | $\backslash$ quant |
| 23 | Intensification | $\backslash$ intens |
| 24 | Comparison | $\backslash$ compar |
| 25 | Correlation | $\backslash$ correl |
| 26 | Substantivation | $\backslash$ substantiv |
| 27 | Nominalization | $\backslash$ nomin |
| 28 | Quotation | $\backslash$ quot |
| 29 | Interrogation | $\backslash$ interr |
| 30 | Imperativization | $\backslash$ imper |

Table 2. GNS Operations
egories). In this section we introduce the forty axioms related to the above thirty syntactic operations. Some examples of their functioning are also provided.

We start by providing the first five axioms of our model according to which the operations of conjugation, determination and predication are applied.

## Axiom 1: Conjugation

$$
\forall x y z(\backslash \operatorname{conjug}(x, y)=z \wedge y: \text { Conjugative } \rightarrow z: \text { Verbt })
$$

The type Conjugative refers to the conjugation parameters that in English, Spanish and Italian are given by Tense, Mood, Person, Number. Examples of conjugative grammemes are: PRES (present), PAST (past), FUT (future), PPG (present-progressive), IND (indicative), SUB (subjunctive), PART (participle), I (first), II (second), III (third), SING (singular), PLUR (plural), NEC (necessity), PERF (perfect), IMP (imperfect), MASC (masculine), FEM (feminine), NEUT (neuter).

## Axiom 2: Determination

$$
\forall x y z(\backslash \operatorname{determ}(x, y)=z \wedge y: \text { Determinative } \rightarrow z: \text { Subst })
$$

The basic determinative grammemes used in determination DEF (definite/near) and INDEF (indefinite/far) in many languages need to be accompanied with grammemes of number (SING (singular), PLUR (plural)) and of gender (MASC (masculine), FEM (feminine), NEUT (neuter)). In other languages the set of determinatives could be richer and related to an intrinsic mechanism of noun classification.

Predication is the basis of the grammatical schema that provides a sentence by combining a 'subject' and a 'predicate'. The following axioms show how to apply this operation.

## Axiom 3: Predication

$$
\forall x(x: \operatorname{Prop} \rightarrow \exists u, v(u: \text { Subst } \wedge v: \operatorname{Verbt} \wedge x=\backslash \operatorname{pred}(u, v)))
$$

## Axiom 4: Predication

$\forall u v w z((u: S u b s t \wedge v: V e r b t \wedge w: V e r b \wedge v \neq \lambda \wedge z:$ Conjugative $\wedge v=$ $\backslash \operatorname{conjug}(w, z) \wedge \operatorname{Agree}(u, v)) \rightarrow \backslash \operatorname{pred}(u, v): \operatorname{Prop})$

## Axiom 5: Predication

$\forall x u v(x: \operatorname{Prop} \wedge x=\backslash \operatorname{pred}(u, v) \rightarrow v=\operatorname{PREDICATE}(x) \wedge u=\operatorname{SUBJ}(x)$
Knowing how conjugation, determination and predication work, we can consider the sentence: 'The dog barks'. It can be formalized as follows (by $D E F$ we mean the definite determiner):

$$
\backslash \operatorname{pred}(\backslash \operatorname{determ}(\operatorname{dog}, \mathrm{DEF}), \backslash \operatorname{conjug}(\mathrm{bark}, \text { III, SING, PRES)})
$$

The deep linguistic level of abstract syntax representations is easily understood if we consider the Italian and Spanish translations of the above sentence: 'Il cane abbaia' and 'El perro ladra', respectively, whose abstract syntax representations are the following:
$\backslash \operatorname{pred}(\backslash \operatorname{determ}($ cane, DEF $), \backslash \operatorname{conjug}($ abbaiare, III, SING, PRES))
$\backslash \operatorname{pred}(\backslash \operatorname{determ}($ perro, DEF), $\backslash \operatorname{conjug}$ (ladrar, III, SING, PRES))
where we can see that these two representations can be obtained from the English one by replacing the lexical items $d o g$ and bark with the corresponding lexical items in Italian and Spanish.

Copulation is an operation that transforms a Noun and AdNoun, or a Subst into a Verb. Axiom 6 provides the requirements to apply this operation.

## Axiom 6: Copulation

$$
\forall x y((x: \text { Noun } \vee x: \text { AdNoun } \vee x: \text { Subst }) \wedge y: \text { Copulative } \rightarrow \backslash \operatorname{copul}(x, y):
$$

Copulative grammemes are BE, BECOME, SEEM. . . Examples of copulative constructions are the following:

- 'was a lawyer': \conj(III, SING, PAST ( $\backslash$ copul(lawyer, BE))
- 'became president': \conj(III, SING, PAST (\copul(president, BECOME))

Modification $(\backslash \operatorname{modif})$ is the combination of a phrase of category $A d C a t$ with a phrase of a category Cat. Modification introduces the notions of kernel and modifier. Axioms 7 and 8 formalize this operation:

## Axiom 7: Modification

$$
\forall x y(x: C a t \wedge y: \operatorname{AdCat} \wedge \operatorname{Agree}(x, y) \rightarrow \backslash \operatorname{modif}(x, y): C a t)
$$

## Axiom 8: Modification

$$
\begin{gathered}
\forall x(\exists u v: u: C a t \wedge v: A d C a t \wedge x=\backslash \operatorname{modif}(u, v) \rightarrow \operatorname{KER}(x)= \\
u \wedge \operatorname{MODIFIER(x)=v)}
\end{gathered}
$$

The following examples show how modification works in GNS:

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- 'young artist': $\backslash \operatorname{modif}($ artist, young)
- 'to walk slowly': \modif(walk, \adverb(slow))

Apposition ( $\backslash$ appos) is an expansion of a Subst with a Noun or an $A d N o u n$. The noun or $A d N o u n$ does not add any further element into the determination of the substantive, but it gives only additional descriptive aspects of what is identified by the substantive. Axiom 9 accounts for this operation.

## Axiom 9: Apposition

$$
\forall x y(x: S u b s t \wedge y: N o u n \vee y: \operatorname{AdNoun} \rightarrow \backslash \operatorname{appos}(x, y): \text { Subst })
$$

An example of apposition is the following:

- 'Rome, the capital of Italy': \appos(Rome, capital of Italy)

Two basic axioms that refer to agreement features are the following. Note that axiom 10 axiomatizes the requirements for the relation Agree that guarantees the presence of some common features between the subject and the predicate of a predication. Axiom 11 concerns some commutativity requirements between the kernel of a modification and the agreement features of a verb or a substantive. We indicate by Feature a typical feature (e.g. gender, number, person, ...)

## Axiom 10: Agreement

$$
\forall x y z(\backslash \operatorname{pred}(x, y)=z \rightarrow(x: \operatorname{Subst} \wedge y: \operatorname{Verb} \wedge \operatorname{Agree}(x, y)))
$$

## Axiom 11: Agreement

$$
\forall x y(x: C a t \wedge y: A d C a t \rightarrow F e a t u r e(x)=\text { Feature }(\backslash \operatorname{modif}(x, y)))
$$

Axioms 12, 13, 14 and 15 show how complementation is applied. Complementation ( $\backslash$ complem) is an operation in which a Noun or a Verb is expanded either with a Subst or with a Noun. This means that we can identify four possible types of complementation:

1. $\backslash$ complem $_{V S}$ (a verb with a substantive);
2. complem $_{N S}$ (a noun with a substantive);
3. $\backslash$ complem ${ }_{V N}$ (a verb with a noun);
4. $\backslash$ complem $m_{N N}$ (a noun with a noun).

These four types of complementation are formalized in the following axioms:

## Axiom 12: Complementation

$$
\begin{gathered}
\forall x y z\left(x: V e r b \wedge y: S u b s t \wedge z: \text { Complementative } \rightarrow \backslash \operatorname{complem}_{V S}(x, z, y):\right. \\
\text { Verb })
\end{gathered}
$$

## Axiom 13: Complementation

$$
\begin{gathered}
\forall x y z\left(x: \text { Noun } \wedge y: S u b s t \wedge z: \text { Complementative } \rightarrow \backslash \operatorname{complem}_{N S}(x, z, y):\right. \\
\text { Noun })
\end{gathered}
$$

## Axiom 14: Complementation

$$
\begin{gathered}
\forall x y z\left(x: \operatorname{Verb} \wedge y: \text { Noun } \wedge z: \text { Complementative } \rightarrow \backslash \operatorname{complem}_{V N}(x, z, y):\right. \\
\text { Verb })
\end{gathered}
$$

## Axiom 15: Complementation

```
\(\forall x y z\left(x:\right.\) Noun \(\wedge y:\) Noun \(\wedge z:\) Complementative \(\rightarrow \operatorname{complem}_{N N}(x, z, y):\)
Noun)
```

Complementative are grammemes that identify the roles of complementations. They correspond to the cases typical of many languages (e.g. Latin, Greek, Russian, ...) and can be represented with specific grammemes related to a few main functionalities: AGT (agent), INST (instrument); GOAL (goal), MANN (manner), MATT (matter), OWN (owner), REF (referent), REC (receiver), UNI (union), CONT (content), PURP (purpose), and CAUS (cause).

Examples of complementation are the following:

- 'Jane eats the cracker': \} complem _ { V S } ( eats, the cracker).
- 'cup of tea': \complem ${ }_{N N}($ cup, CONT, tea).

Locative phrases are a special type of modifiers used to indicate a spatiotemporal localization. Axiom 16 accounts for this syntactic phenomenon which we call localization ( $\backslash l o c$ ):

## Axiom 16: Localization

$$
\forall x y(x: \text { Subst } \wedge y: \text { Locative } \rightarrow \backslash \operatorname{loc}(x, y): \text { AdCat })
$$

Locative grammemes can be IN, OUT, TO, FROM, AROUND, ACROSS, BETWEEN, NEAR, OVER, TOWARDS, BEFORE, AFTER, etc.
An example of a locative phrase is:

- '(Jane) found the turtle under the table': $\backslash \operatorname{modif}\left(\backslash \operatorname{complem}_{V S}\right.$ (found, the turtle), $\backslash \operatorname{loc}($ the table, UNDER))

Coordination ( $\backslash$ coord) allows us to connect two (or more) categories of the same type by providing a category of that type. If two elements are of the same category, then they can be joined together and the resulting unit is of the same type. Any category can be coordinated (by means of a coordinative grammeme: AND (addition), OR (alternation), BUT (variation)) with a category of the same type and the result will be an item of the same category. Axiom 17 formalizes this idea.

## Axiom 17: Coordination

$$
\forall x y z(x: C a t \wedge y: C a t \wedge z: C o o r d i n a t i v e \rightarrow \backslash \operatorname{coord}(x z y): C a t)
$$

Consider the sentence 'We saw many students of chemistry and doctors of medicine', where the elements in italics are coordinated. In terms of GNS, coordinated categories can be formalized as follows:
$\backslash \operatorname{coord}\left(\left(\backslash \operatorname{complem}_{N N}(\right.\right.$ students, chemistry $\left.)\right)$, AND, $\left(\backslash \operatorname{complem}_{N N}\right.$ (doctors, medicine)))

Subordination connects elements that have a different grammatical 'status', one of which is subordinate to or dependent on the other. Subordinative grammemes are: BECAUSE (reason), IF (condition), WHERE (place), WHEN, AFTER, BEFORE (time), WHILE, UNTIL (duration), THOUGH (concession), FOR, TO (purpose). Axioms 18 and 19 give a formal account in GNS of this operation.

## Axiom 18: Subordination

$$
\forall x y z(x: C a t \wedge y: C a t \wedge z: \text { Subordinative } \rightarrow \backslash \operatorname{subord}(x z y): C a t)
$$

## Axiom 19: Subordination

$\forall x y z(x: C a t \wedge y: C a t \wedge z:$ Subordinative $\rightarrow \backslash \operatorname{subord}(x z y):$ Cat $\wedge x=$ $\operatorname{Ker}(\backslash \operatorname{subord}(x z y)))$

Examples of the subordination of various classes of categories are: 'poor but happy', 'late though not too late', 'enemy for joke', 'going to see'... The reader can easily check how to formalize these examples by using the above axiom.

Negation $(\backslash n e g)$ is a syntactic operation that can be applied to any Cat and the result will be an item of the same Cat. Axiom 20 formalizes this operation.

## Axiom 20: Negation

$$
\forall x(x: C a t \rightarrow \backslash n e g(x): C a t)
$$

Examples of the application of this operation in the GNS framework are:

- '(I do) not agree': $\backslash n e g$ (agree)
- '(It is) not soon': $\backslash n e g($ soon $)$

Passivation ( $\backslash$ passiv) is an operation that like negation does not change the syntactic category of the element it is applied to. It is based on the fact that the order of the subject and object can be reversed in a predication with transitive verbs. In this case, if P is a proposition to which passivation can be applied, then $\backslash \operatorname{passiv}(\mathrm{P})$ is its passive form. The actual means to calculate the value of $\backslash \operatorname{passiv}(\mathrm{P})$ is a matter of concrete syntax. Formally:

## Axiom 21: Passivation

$$
\forall x(x: \operatorname{Prop} \rightarrow \backslash \operatorname{passiv}(x): \text { Prop })
$$

Axioms 22, 23 and 24 formalize cases of transcategorization. Adcategorization stands for the operations of adjectivation, adverbialization and adproposition considered in the following axioms.

Adjectivation $(\backslash a d j)$ takes as its input a Noun, a Subst, a Verb or a Verbt and provides an $A d N o u n$ as shown in axiom 22.

## Axiom 22: Adjectivation

$\forall x y((x: N o u n \vee x: S u b s t \vee x: V e r b \vee x: V e r b t) \wedge \backslash \operatorname{adj}(x): A d N o u n)$

Adverbialization ( $\backslash$ adverb) takes as an input an AdNoun and provides an AdVerb.

## Axiom 23: Adverbialization

$$
\forall x(x: \operatorname{AdNoun} \rightarrow \backslash \operatorname{adverb}(x): A d V e r b)
$$

And finally, adproposition ( $\backslash$ adprop) yields an adproposition element from either an AdNoun or a Subst.

## Axiom 24: Adproposition

$$
\forall x(x: \text { AdNoun } \vee x: \text { Subst } \rightarrow \backslash \operatorname{adprop}(x): \text { AdProp })
$$

Examples of the above operations are the following:

- 'milk cup': $\backslash \operatorname{modif}(\mathrm{cup}, \backslash \operatorname{adj}(\mathrm{milk}))$
- 'slowly': \adverb(slow)
- 'This way, (Doris feeds her guppies)': \adprop(this way)

Relativization ( $\backslash$ relat) is the syntactic operation that transforms a Prop into a Noun. In order to describe this construction, we assume a special grammeme REL that has the category of Subst (Actually, a finite number of different relative grammemes REL1, REL2, REL3, ..., could be necessary in certain cases). If proposition P is 'REL was on the table'. In this case, $\backslash$ relat ( P ) is a noun, so the following syntactic term:
'I am looking for the $\backslash$ modif (pen, $\backslash \operatorname{adj} \backslash$ relat $(P))$ ' represents the statement: 'I am looking for the pen that was on the table.'

Axiom 25 formalizes relativization, where Prop(REL) are the propositions constructed by the special substantive REL.

## Axiom 25: Relativization

$$
\forall x(x: \operatorname{Prop}(\mathrm{REL}) \rightarrow \backslash \operatorname{relat}(x): \text { Noun })
$$

A syntactic operation formally described by the element of Prop (REL) is the one that underlies the so-called phenomenon of anacoluthon, a construction present in ancient languages and in colloquial forms of modern languages. We introduce the operation of anacoluthon ( $\backslash$ anacoluth) defined in axiom 26.

## Axiom 26: Anacoluthon

$$
\forall x(x: \operatorname{Prop}(\mathrm{REL}) \rightarrow \backslash \text { anacoluth }(x): \text { Verb })
$$

An example of anacoluthon formalized in GNS terms is the following:
'John, I told him to go away.': \pred(John, \conjug(\anacoluth(I told REL to go away), III, SING))

Axioms 27 and 28 formalize the operations of numeralization and ordinalization, respectively. Numeralization ( $\backslash$ num) takes as its arguments a cardinal number (two, three, ten, two hundred, three thousand, a dozen, hundreds) and a noun and provides a noun. However, this operation is only defined for some nouns, which are called countable nouns. Ordinalization ( $\backslash$ ord) takes as its arguments an ordinal number (first, second, last, next...) and a noun and provides a noun.

## Axiom 27: Numeralization

$$
\forall x y(x: \text { Noun } \wedge y: \text { Cardinal } \rightarrow \backslash \operatorname{num}(x, y): \text { Noun })
$$

## Axiom 28: Ordinalization

$$
\forall x y(x: \operatorname{Noun} \wedge y: \text { Ordinal } \rightarrow \backslash \operatorname{ord}(x, y): \text { Noun })
$$

The following two phrases are examples of these operations:

- 'Three friends': $\backslash$ num(friend, 3))
- 'The first friend': \determ( $\backslash$ ord(friend, 1), DEF)

Deictification ( $\backslash$ deixis) is an operation that takes as its arguments some grammemes (conjugative, determinative: PRES, PAST, I, II, III, DEF/INDEF, NEAR, FAR) and provides a Subst. Axiom 29 formalizes this idea.

## Axiom 29: Deictification

$$
\forall x(x: \operatorname{Deictic} \rightarrow \backslash \operatorname{deixis}(x): S u b s t)
$$

English examples of deictification are the following:

- $\backslash$ deixis $($ PRES $)=$ 'Now'
- $\backslash \operatorname{deixis}(\mathrm{I}, \mathrm{SING})=$ ' $I$ '
- $\backslash \operatorname{deixis}(\mathrm{NEAR})=$ 'Here’

Anaphorization ( $\backslash$ anaphor) is a construct that assigns a label to syntactic terms (pronoun) which can be used instead of the labeled term itself. Notice that the label (pronoun or whatever) is something that concerns the concrete syntax and not the abstract syntax we are defining here. Taking this into account we define this operation in axiom 30 .

## Axiom 30: Anaphorization

$$
\forall x(x: \text { Cat } \rightarrow \backslash \operatorname{anaphor}(x): \text { Cat })
$$

An example:

- 'On the table there is a pen, it is red': \anaphor(a pen) $=$ 'it'

Quantification ( $\backslash$ quant) can be considered as a special case of determination. However, it presents some subtle logical and semantic features that mean that it has to be identified separately from determination. First of all, quantification can only be applied to countable nouns or their plurals and provides as its result a substantive with a 'collective and distributive' nature. It is formalized in axiom 31.

## Axiom 31: Quantification

$$
\begin{aligned}
& \forall x y(x: \text { Quantifier } \wedge y: \text { Noun }- \text { Countable } \rightarrow \backslash \text { quant }(x, y): \text { Subst }) \\
& \forall x y(x: \text { Quantifier } \wedge y: \text { Subst }- \text { Plural } \rightarrow \backslash \text { quant }(x, y): \text { Subst })
\end{aligned}
$$

Quantifiers are grammemes that express the functionalities of words that classical grammars classify as distributive pronouns: EVERY, EACH, ALL, ANY, FEW, SOME, MOST, MANY, etc. (Some of them are very often referred to as 'partitive' pronouns).

Intensification ( $\backslash$ intens) consists of the use of grammemes that express intensity or non-exact quantity with uncountable nouns. Examples of intensive grammemes are MUCH, A LITTLE, FEW, ALMOST, MORE, LESS, VERY, WHOLE, PARTIAL, ENOUGH, TOO... Axiom 32 provides a formal account of this operation.

## Axiom 32: Intensification

$$
\begin{aligned}
& \forall x y(x: \text { Noun }-U \text { ncountable } \wedge y: \text { intensive } \rightarrow \backslash \text { intens }(x, y): \text { Noun }) \\
& \forall x y(x: \text { AdNoun } \wedge y: \text { intensive } \rightarrow \backslash \text { intens }(x, y): \text { AdNoun })
\end{aligned}
$$

Comparison ( $\backslash$ compar) is related to degree particles such as more-than, less-than, so-as. Comparison can be obtained in three ways: 1) with two terms of a given category Cat and a comparison grammeme (this is, $(<,>,=)$ ); 2) with a Verb or an AdNoun, a comparison grammeme, and a Subst; 3) with a Verb or an AdNoun, a comparison grammeme, and a Prop. This idea is captured in axiom 33.

## Axiom 33: Comparison

```
\(\forall x y z(x: C a t \wedge y: C a t \wedge z:\) Comparative \(\rightarrow \backslash \operatorname{compar}(x, y, z): C a t) \wedge\)
\(\forall x y z T(x: T \wedge \wedge(T=\) Verb \(\vee T=\) AdNoun \() \wedge y:\) Subst \(\wedge z:\) Comparative
\(\rightarrow \backslash \operatorname{compar}(x, y, z): T) \wedge\)
\(\forall x y z T(x: T \wedge(T=V e r b \vee T=\) AdNoun \() \wedge y: \operatorname{Prop} \wedge z:\) Comparative
\(\rightarrow \backslash \operatorname{compar}(x, y, z): T)\)
```

Examples of comparison are the following:

- \compar(beautiful, good, >): 'more beautiful than good'
- \compar(beautiful, Mary, =): 'as beautiful as Mary'
- $\backslash$ compar(good, the original, $=$ ): 'as good as the original'

Axioms 34 and 35 account for the operation we have called correlation. Correlation ( $\backslash$ correl) is found in constructions where two propositions are related by means of an intensivity degree of a Verb or an AdNoun that occurs in the first proposition (e.g. so-that, such-that).

## Axiom 34: Correlation

$$
\begin{gathered}
\forall p q x u v w y(p: \operatorname{Prop} \wedge q: \operatorname{Prop} \wedge p=u w v \wedge(w: \operatorname{Verb} \vee w: \operatorname{AdNoun}) \wedge w= \\
\backslash \operatorname{intens}(x, y) \rightarrow \operatorname{correl}(p, w, q): \operatorname{Prop})
\end{gathered}
$$

## Axiom 35: Correlation

$$
\begin{gathered}
\forall p q x u v w z y r s t((p: \operatorname{Prop} \wedge p=u w v \wedge(w: \operatorname{Verb} \vee w: \operatorname{AdNoun}) \wedge(q: \\
\operatorname{Prop} \wedge q=z x y \wedge(x: \operatorname{Verb} \vee x: \operatorname{AdNoun}) \wedge t=\backslash \operatorname{compar}(w, r, \mathrm{REL}) \wedge d= \\
\backslash \operatorname{compar}(x, s, \mathrm{REL}) \rightarrow \operatorname{correl}(p, t, q, d): \operatorname{Prop})))
\end{gathered}
$$

Examples of those two types of correlation are the following:

- '(This is) so big that you cannot carry it.': \correl(This is very big, $\backslash i n t e n s($ big, VERY), you cannot take it)
- 'The quicker you are, the more they appreciate your work': \correl(you are quick, $\backslash$ compar (quick, $>$, REL), they appreciate your work, $\backslash$ compar(appreciate, $>$, REL).

Substantivation ( $\backslash$ substantiv) and nominalization ( $\backslash$ nomin) transcategorize some categories into Subst (substantives) and Noun (nouns), respectively. According to a general principle of language, any linguistic entity can become something that can be spoken about; that is, any category can be nominalized and substantivized. We formalize such general principles with axioms 36 and 37 , which account for substantivation and nominalization ( $\backslash$ nomin), respectively.

## Axiom 36: Substantivation

$$
\forall x(x: \text { Prop } \vee x: \text { Noun } \vee x: \text { AdNoun } \rightarrow \backslash \operatorname{substantiv}(x): \text { Subst })
$$

## Axiom 37: Nominalization

$$
\forall x(x: V e r b \vee x: A d N o u n \rightarrow \backslash \operatorname{nomin}(x): \text { Noun })
$$

Examples of the two operations above are:

- 'eating': \substantiv( $\backslash$ nomin(to eat))
- 'emptiness': \nomin(empty)

A particular case of substantivation is quotation ( $\backslash q u o t$ ) which transforms any sequence of letters into a Subst. This particular case of substantivation is formalized in axiom 38.

## Axiom 38: Quotation

$$
\forall x(x: C a t \rightarrow \backslash q u o t(x): \text { Subst })
$$

Consider the following sentence:

## "He said: 'Tomorrow I will finish the paper"

Here a full proposition works as a substantive. So,
'Tomorrow I will finish the paper': \quot (Tomorrow I will finish the paper)

The last two operations in the GNS framework are interrogation ( $\backslash$ interr ) and imperativation ( $\backslash$ imperat). Those general syntactic operations transform any descriptive proposition into the corresponding proposition with a communication modality which is interrogative, or imperative respectively. Axioms 39 and 40 account for them.

## Axiom 39: Interrogation

$$
\forall x(x: \operatorname{Prop} \rightarrow \backslash \operatorname{interrog}(x): \text { Prop })
$$

## Axiom 40: Imperativation

$$
\forall x(x: \text { Prop } \rightarrow \backslash \operatorname{imperat}(x): \text { Prop })
$$

## 4 An Example

In this section we provide an example of how GNS works. First, we analyze the following English sentence, and then we show that the formal representation of this English sentence - by just translating basic words- can be used to generate equivalent Italian and Spanish sentences.
"The two children ran towards the river with their hands raised till they reached the bank"

The above English sentence is analyzed in the following way (we use a \# to indicate the line number):

1. hand
2. to raise
3. child
4. river
5. to run
6. to reach
7. bank
8. $\backslash \operatorname{conjug}(\# 2$, PART,$~ P A S T)=$ raised
9. $\backslash \operatorname{modif}(\# 1, \backslash \operatorname{adj}(\# 8))=$ hand raised
10. $\backslash \operatorname{determ}(\# 9$, DEF, PLUR $)=$ their hands raised
11. $\backslash \operatorname{num}(\# 3,2)=$ two children
12. $\backslash \operatorname{determ}(\# 11, \mathrm{DEF})=$ the two children
13. $\backslash \operatorname{conjug}(\# 5$, III, PLUR, PAST, IMP $)=r a n$
14. $\backslash \operatorname{determ}(\# 4, \mathrm{DEF})=$ the river
15. $\backslash \operatorname{loc}(\# 14$, TOWARDS $)=$ towards the river
16. $\backslash$ complem $(\# 13, \# 15)=$ ran towards the river
17. $\backslash \operatorname{complem}(\# 16$, MANN, $\# 10)=$ ran towards the river with their hands raised
18. $\backslash \operatorname{pred}(\# 12, \# 17)=$ the two children ran towards the river with their hands raised
19. $\backslash \operatorname{conjug}(\# 6$, III, PLUR, PAST, PERF $)=$ reached
20. $\backslash \operatorname{determ}(\# 7, \mathrm{DEF})=$ the bank
21. $\backslash$ complem $_{V S}(\# 19, \# 20)=$ reached the bank
22. $\backslash \operatorname{pred}(\backslash$ deixis(III, PLUR),$\# 21)=$ they reached the bank
23. $\backslash$ subord $(\# 18$, UNTIL, $\# 22)=$ The two children ran towards the river with their hands raised till they reached the bank.

In the end we have the basic elements and operations shown in Table 3.

| BASIC ELEMENTS OPERATIONS |  |
| :--- | :--- |
| child | $\backslash$ subord |
| to run | $\backslash$ pred |
| river | $\backslash$ determ |
| hand | $\backslash$ num |
| to raise | $\backslash$ complem |
| to arrive | $\backslash$ conjug |
| bank | $\backslash$ deixis |
|  | $\backslash$ loc |
|  | $\backslash$ modif |

Table 3: Basic Elements and Operations.
If we now take the above GNS formalization and just translate the so-called basic elements into Italian or Spanish in Table 4 below.

| English | Italian | Spanish |
| :--- | :--- | :--- |
| child | bambino | niño |
| to run | correre | correr |
| river | fiume | río |
| hand | mano | mano |
| to raise | alzare | levantar |
| to reach | raggiungere | alcanzar |
| bank | riva | orilla |

Table 4: Translation of Basic Elements
we obtain the following two sentences:

- I due bambini correvano verso il fiume con le mani alzate finché non raggiunsero la riva.
- Los dos niños corrieron hacia el río con las manos levantadas hasta que alcanzaron la orilla.


## 5 Conclusions

In this paper we have presented a formal theory of general explicative power: General Natural Syntax. The aim behind this framework is to reduce syntactic constructions to a few principles related to their semantic functions but defined independently of semantics. In this regard, it can be related to the socalled abstract syntax of programming languages. In fact, GNS establishes the requirements for the construction of terms and the rules for assigning them syntactic categories. It does not deal with the final surface syntactic forms of concrete syntax. It does not try to cope with the particular morphological features of the final linguistic form of the expression resulting from the application of the syntactic operation, but it is interested in a deeper syntactic level: GNS aims to define an abstract syntax for natural languages.

Within GNS, many classical notions can be stated formally, in terms of basic categories and basic abstract syntactic operations. GNS has the advantage of simplicity. With a very small number of simple ingredients, GNS can account for general syntactic constructions in any natural language. GNS generalizes and formalizes syntactic phenomena present in every natural language and incorporates them into a formal theory of general explicative power. Although GNS is very recent and still requires further research, if its features
are taken into account we feel that it may be suitable in the fields of natural language processing, machine translation and linguistics.

Our future work will focus on two main aspects: 1) basic lexicon and grammemes; and 2) formal and computational work. Regarding the former, we are interested in establishing a basic lexicon and an exact number of grammemes. Regarding the latter, it would be very useful to define an interface to generate GNS formula and to develop a system that can generate sentences from the GNS formula. Besides those two topics, it would be very interesting to go deeper into the possible applications of the framework.

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[^0]:    ${ }^{3}$ Note that we are not using 'grammeme' here with its usual meaning in linguistic morphology.

